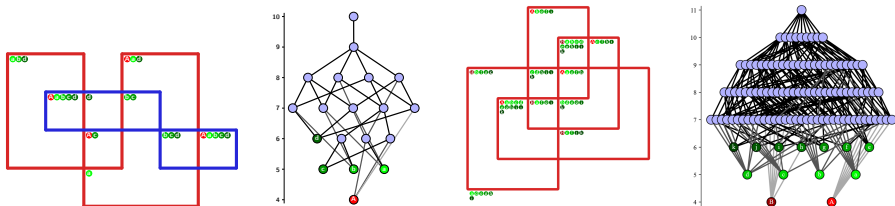


The complexity of pinning loops in the plane

Christopher-Lloyd SIMON and Ben Stucky

The Pennsylvania State University

Geometry Lunch Seminar, 2024-10-30



- 1 Why study combinatorics and complexity of plane loops ?
- 2 Multiloops and their pinning sets
- 3 From pinning-loops to hypergraphs via immersed discs
- 4 Proof of the MoBiDisc theorem
- 5 Experimentations and further directions of research

Complexity of topological invariants

Recent results on the computational complexity of link invariants:

bounding the **crossing number** of a knot or link is NP-hard

bounding the **unknotting number** of a knot or link is NP-hard

bounding the **genus** of a knot in a general 3-manifold is NP-hard

bounding the **genus** of a knot in the sphere is in co-NP

computing a **finite type invariant** of degree d is in $O(n^d)$

Handwritten notes by Gauss:

Die Zahl der Ueberschneidungen zweier Linien ist die positive oder negative Zahl der Ueberschneidungen zweier Linien, die sich nicht selbst schneiden.

Die Zahl der Ueberschneidungen zweier Linien ist die positive oder negative Zahl der Ueberschneidungen zweier Linien, die sich nicht selbst schneiden.

Die Zahl der Ueberschneidungen zweier Linien ist die positive oder negative Zahl der Ueberschneidungen zweier Linien, die sich nicht selbst schneiden.

Die Zahl der Ueberschneidungen zweier Linien ist die positive oder negative Zahl der Ueberschneidungen zweier Linien, die sich nicht selbst schneiden.

DER ELECTRODYNAMIK. 605

[4.]

Von der Geometria Situs, die LEONARDI ahlte und in die nur einem Paar Geometers (EUCLID und VALENTINUS) einen schwachen Blick zu thun vergabten war, wissen und haben wir nach anderthalbhundert Jahren noch nicht viel mehr als nichts.

Eine Hauptaufgabe aus dem Grenzgebiet der Geometria Situs und der Geometria Magnitudinis wird die sein, die Umschlingungen zweier geschlossener oder unendlicher Linien zu zählen.

Es seien die Coordinaten eines unbestimmten Punktes der ersten Linie x, y, z ; der zweiten x', y', z' und

$$\iint \frac{(x-x')(y-y') + (z-z')(y-y') - (x-x')(z-z') + (y-y')(z-z') - (x-x')(z-z') - (y-y')(z-z')}{[(x-x')^2 + (y-y')^2 + (z-z')^2]^{3/2}} dx dy dz dx' dy' dz'$$

dann ist dies Integral durch beide Linien ausgedehnt

= 4 NT

und so die Anzahl der Umschlingungen.

Der Werth ist gegenseitig, d. i. er bleibt derselbe, wenn beide Linien gegen einander umgetauscht werden. 1833. Jun. 22.

Gauss discovers linking numbers.

Everything is contained in plane loops

Arnol'd [Arn00] in the Preface to Lecture 1:

A method for studying infinite dimensional configuration space is to compute the relative cohomology of their stratifications by singular loci. This is how Vassiliev introduced finite type invariants of knots $\mathbb{S}^1 \hookrightarrow \mathbb{S}^3$. Arnol'd proposes an analogous study for immersions $\mathbb{S}^1 \looparrowright \mathbb{S}^2$, but...

The combinatorics of plane curves seems to be far more complicated than that of knot theory (which might be considered as a simplified “commutative” version of the combinatorics of plane curves and which is probably embedded in plane curves theory).

Survey [Poé95] on extending immersions $\mathbb{S}^m \looparrowright \mathbb{S}^n$ to $\mathbb{D}^{m+1} \looparrowright \mathbb{S}^n$

Thom-Smale-Hirsch: solved using homotopy groups and obstruction theory when $m - n > 1$. Note: homotopy groups are commutative in degrees > 1 .

Why self-intersection and pinning sets of multiloops ?

A (whimsical) physical motivation: Neumann-Coto [NC01]

In a surface \mathbb{F} , an immersion $\gamma: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{F}$ is isotopic to a collection of **shortest geodesics** for a complete Riemannian metric on \mathbb{F} if and only if it is **taut** (realizes the minimal number of double-points in its homotopy class).

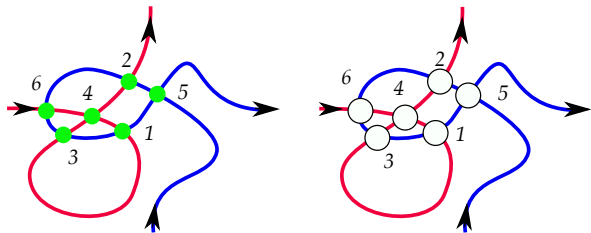
Where do punctures come into the picture ?

For multiloops in the sphere, we will study the relation between self-intersection numbers and punctures to define pinning numbers. Applications to the bounded cohomology of braid groups and $\text{Homeo}(\mathbb{S}^2)$.

Using multiloops as a computation scheme

Vaughan Jones: planar algebras

The algebra of tangles and multiloops in punctured spheres encapsulate efficiently the essence (and complexity) of many computational schemes.



Computing with planar algebras.

- 1 Why study combinatorics and complexity of plane loops ?
- 2 Multiloops and their pinning sets**
- 3 From pinning-loops to hypergraphs via immersed discs
- 4 Proof of the MoBiDisc theorem
- 5 Experimentations and further directions of research

Multiloops in punctured spheres

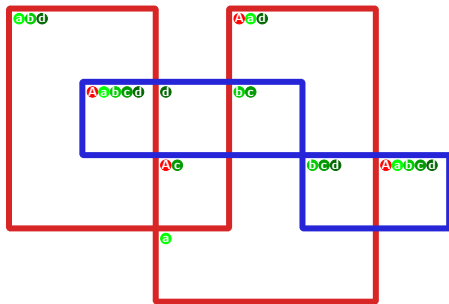
Disjoint union of circles $\sqcup_1^s \mathbb{S}^1$. Punctured sphere $\mathbb{S}_P^2 = \mathbb{S}^2 \setminus P$.

A **multiloop** with s strands is a generic immersion $\gamma: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$ up to orientation preserving diffeomorphisms of source and the target.

Here **generic** means all multiple points are transverse double-points.

The **number of double-points** is denote by $\#\gamma$.

The **regions** R are the connected components of $\mathbb{S}^2 \setminus \text{im}(\gamma)$.



The multiloop 10_{16}^2 has $\#\gamma = 8$. Various sets of punctures.

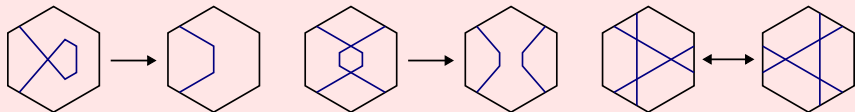
Multicurves, self-intersection, taut multiloops

A **multicurve** is a homotopy class of multiloops.

The **self-intersection** $si(\gamma)$ of a multicurve is the minimal number of double-points among its generic representatives.

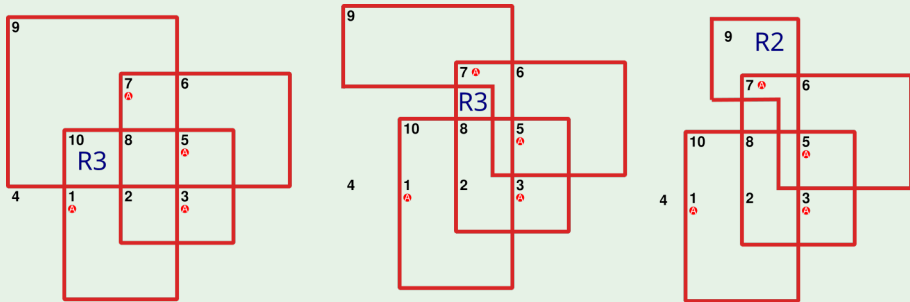
A multiloop γ is **taut** when $\#\gamma = si(\gamma)$, namely when it has the minimal number of double-points in its homotopy class.

If two multiloops in \mathbb{S}_P^2 are homotopic, then they are related by a sequence of isotopies and Reidemeister moves $R1, R2, R3$.



How to know when a multiloop is taut ?

Try combinations of $R1, R2, R3$ until finding γ' with $\#\gamma' < \#\gamma \dots$



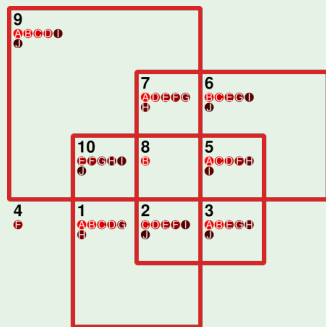
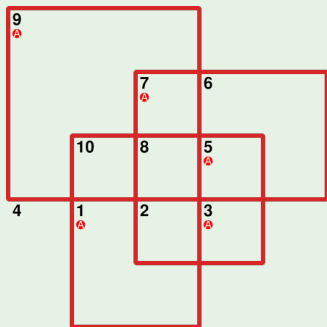
Can we order and bound the sequences of R_k -moves ? We won't insist...

Can we construct an algorithm computing $si(\gamma)$ to compare with $\#\gamma$?

How to know when a multiloop is taut ?

Pinning sets

Which puncture sets $P \subset R$ will yield a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$ that is taut ?



The set $P = R$ of all regions works. We can try deleting one pin at a time...

Pinning sets, pinning number, pinning ideal

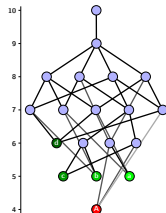
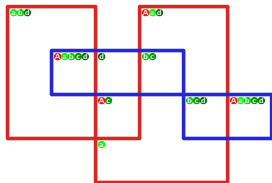
Consider a multiloop $\gamma: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{S}^2$ with regions R .

A subset of regions $P \subset R$ is **pinning** when $\gamma_P: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$ is taut.

The pinning sets of γ form its **pinning ideal**: a subposet $\mathcal{PI} \subset \mathcal{P}(R)$ absorbant under union and containing R .

The **pinning number** $\varpi(\gamma)$ is the minimum cardinal of pinning sets.

A pinning set **optimal** if it has the minimum cardinal, in particular it must be **minimal** (with respect to inclusion).



The multiloop 10_{16}^2 and its pinning semi-lattice obtained from unions of minimal pinning sets (in red or green), together with the whole set R .

Main questions:

Given a multiloop $\gamma: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{S}^2$, how (efficiently) can we:

Construct minimal or optimal pinning sets?

Find which regions belong to every pinning set?

Compute the pinning number $\varpi(\gamma)$?

What can be the shape of the pinning ideal of a (multi)loop?

What are its statistics for certain random models?



The loop 9_5^1 has 2 optimal pinning sets, and 3 other minimal pinning sets.

Overview of our main complexity results

Computing si and check optimal solutions in P-steps

We construct a polynomial algorithm which given a multiloop

$\gamma: \sqcup_1^s \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$ computes its self intersection-number $si(\gamma)$.

Thus in P-steps we can: **check if a given set** $P \subset R$ is pinning, if so minimal and optimal; and we can **find a minimal** pinning set.

We implemented it to build an **online catalog** of pinning data for multiloops.

The Pinning number is NPC:

The following LoopPiNum decision problem is NPC:

Input: A loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}^2$ and an integer $k \in \mathbb{N}$.

Output: Is $\varpi(\gamma) \leq k$?

Proof: we will construct two reductions:

NP-easy: reduction from LoopPiNum to hypergraph-vertex-cover.

NP-hard: reduction from graph-vertex-cover to LoopPiNum.

- 1 Why study combinatorics and complexity of plane loops ?
- 2 Multiloops and their pinning sets
- 3 From pinning-loops to hypergraphs via immersed discs**
- 4 Proof of the MoBiDisc theorem
- 5 Experimentations and further directions of research

Singular monorbignons

Singular monorbignons of a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$:

A **singular monogon** is a non-trivial closed interval $I \subset \mathbb{S}^1$ such that $\gamma(\partial I) = \{x\}$ and $\gamma(I) \subset \mathbb{S}_P^2$ is null-homotopic.

A **singular bigon** is a disjoint union of non-trivial closed intervals $I \sqcup J \subset \mathbb{S}^1$ such that $\gamma(\partial I) = \{x, y\} = \gamma(\partial J)$ and $\gamma(I \sqcup J) \subset \mathbb{S}_P^2$ is null-homotopic. It is **embedded** when γ is injective on the interior of $K = I$ or $K = I \sqcup J$.

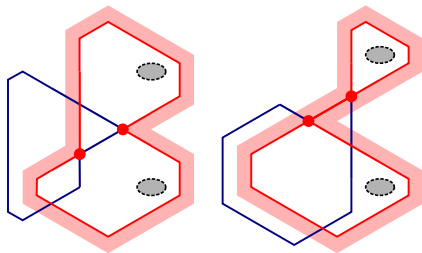
	Regional	Embedded (not regional)	Immersed (not embedded)	Singular (not immersed)	Weak (nonsingular)
1gon					
2gon					

Singular monorbignons

Singular monorbignons of a loop $\gamma: \mathbb{S} \looparrowright \mathbb{S}_p^2$:

A **singular monogon** is a non-trivial closed interval $I \subset \mathbb{S}^1$ such that $\gamma(\partial I) = \{x\}$ and $\gamma(I) \subset \mathbb{S}_p^2$ is null-homotopic.

A **singular bigon** is a disjoint union of non-trivial closed intervals $I \sqcup J \subset \mathbb{S}^1$ such that $\gamma(\partial I) = \{x, y\} = \gamma(\partial J)$ and $\gamma(I \sqcup J) \subset \mathbb{S}_p^2$ is null-homotopic. It is **embedded** when γ is injective on the interior of $K = I$ or $K = I \sqcup J$.



A singular (embedded) bigon $R3$ -moves into a non-singular (weak) bigon.

Hass-Scott characterise tautness with singular monorbigons

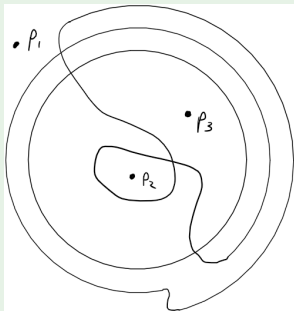
Hass-Scott [HS85, Theorems 2.7 and 4.7]

2.7: If $\#\gamma > 0$ but $\text{si}(\gamma) = 0$, then γ has an embedded monorbigon.

4.2: If $\#\gamma > \text{si}(\gamma)$, then γ has a singular monorbigon.

Just for loops: no generalisation to multiloops

A multiloop $\gamma: \mathbb{S}^1 \sqcup \mathbb{S}^1 \looparrowright \mathbb{S}^2 \setminus \{p_1, p_2, p_3\}$ which is *not taut*, but with *no singular monorbigons* (generalised to multiloop):

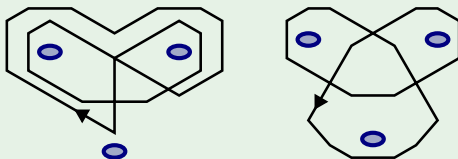


Winding numbers help to pin loops

Corollary: winding numbers yield sufficient conditions for pinning sets

Consider a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}^2$ and a set of regions $P \subset R$. If P is not pinning then there exists a singular monorbigon α of γ such that for every pair of regions $o, \infty \in P$ we have $\text{lk}(\alpha, [o] - [\infty]) = 0$.

The converse is false: the commutator loop in $\mathbb{S}^2 \setminus \{0, 1, \infty\}$



Strategy

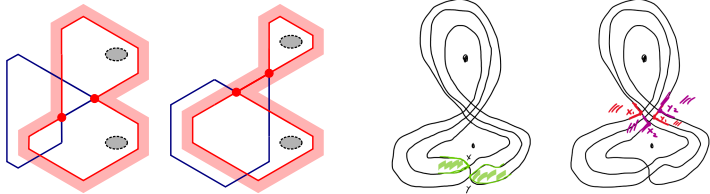
To characterise pinning sets of loops in terms of singular monorbigons and linking numbers, we will restrict to the class of **immersed** monorbigons... These will define the obvious collections of regions to pin: the **mobidiscs**.

Immersed monorbignons and the Mobidisc theorem

Definition: For a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$,

A singular monorbignon $K \subset \mathbb{S}^1$ is **immersed** when there exists $\iota: \mathbb{D} \looparrowright \mathbb{S}_P^2$ such that the restriction $\gamma: K \looparrowright \mathbb{S}_P^2$ factors through $\iota: \partial\mathbb{D} \looparrowright \mathbb{S}_P^2$.

A **mobidisc** consists of the regions in the image $\iota(\mathbb{D})$ of such an immersion. Denote by $\text{MoB}(\gamma) \subset \mathcal{P}(R)$ the set of all mobidiscs.



The MoBiDisc Theorem [SS24]

For a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}^2$ and a collection of regions $P \subset R$.

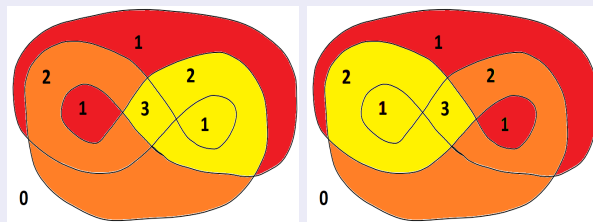
The loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}_P^2$ is taut if and only if it has no immersed monorbignons.

Thus $P \in \mathcal{P}(R)$ is pinning if and only if it intersects every $D \in \text{MoB}(\gamma)$.

Computing MoBiDiscs in polynomial time

Non-unique immersed discs but unique mobidisc

An immersed monorbison may bound *several non-isotopic immersed discs*:



but they all define *the same mobidisc*: its regions are distinguished by $(p, q) \in R \times R \mapsto \text{sign lk}(\gamma(K), [p] - [q]) \in \{0, \pm 1\}$.

Polynomial computation of the set of mobidiscs

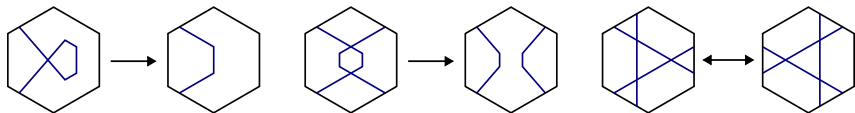
Given a loop $\gamma: \mathbb{S}^1 \looparrowright \mathbb{S}^2$, relying on [Bla67, Fri10] and linking numbers, we can compute $\text{MoB}(\gamma) \subset \mathcal{P}(R)$ time $O(\text{Card}(R)^4)$.

- 1 Why study combinatorics and complexity of plane loops ?
- 2 Multiloops and their pinning sets
- 3 From pinning-loops to hypergraphs via immersed discs
- 4 Proof of the MoBiDisc theorem**
- 5 Experimentations and further directions of research

Outline of the proof

Theorem [HS94]: curve shortening flow

If two loops in \mathbb{S}_p^2 are homotopic, then they are related by a sequence of isotopies and **shortening** Reidemeister moves $R1, R2, R3$.



Lemma: extending immersed monorbignons across R moves

Consider a loop β and a shortening-Reidemeister move to a loop α . If α has an immersed monorbignon then β has an immersed monorbignon.

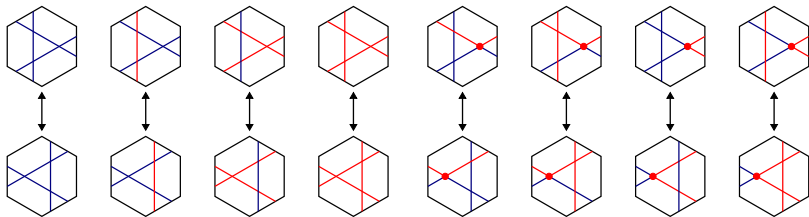
Proof: Chasing mobidiscs accross $R3$ moves by case analysis:

Easy cases: the move restricts to a regular isotopy of the monorbignon.

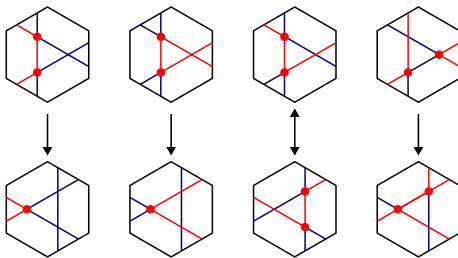
Medium cases : tweak local structure to find immersed monorbignon.

Hard cases: global constructions and arguments to find new mobidisc.

Proof of the Lemma



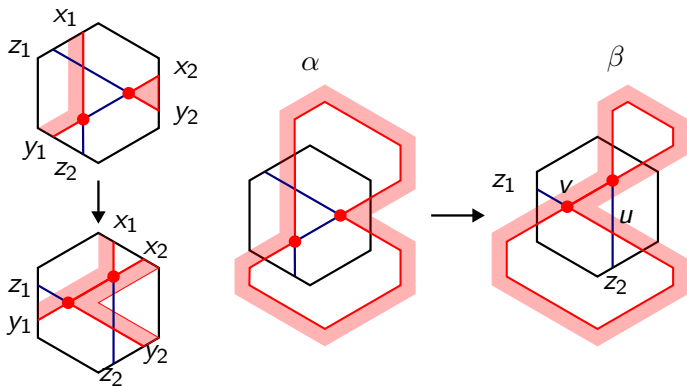
Configurations of type $R3(0, j)$. Configurations of type $R3(1, j)$.



Configurations of type $R3(2, j)$.

Proof of the Lemma: hardest case

Last case: the singular bigon disappears into a weak bigon, so we must construct another singular monorbigon which is immersed.

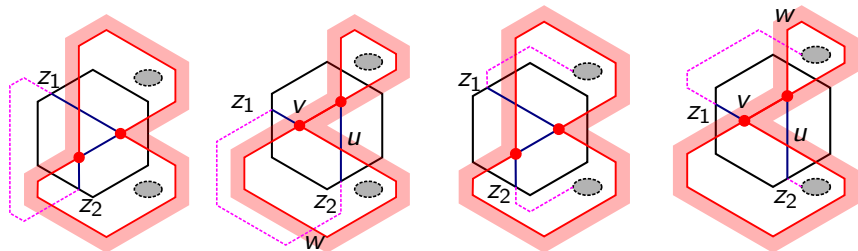


The move $R3(2,2)D$ leads to a weak bigon.

Proof of the Lemma: hardest case

Last case: the singular bigon disappears into a weak bigon, so we must construct another singular monorbigon which is immersed.

Subcase analysis: the construction and proof use orientability of the sphere, Theorems 2.7, & 4.2 of Hass-Scott, and more case analysis.



Finding an immersed monorbigon in cases x and y .

- 1 Why study combinatorics and complexity of plane loops ?
- 2 Multiloops and their pinning sets
- 3 From pinning-loops to hypergraphs via immersed discs
- 4 Proof of the MoBiDisc theorem
- 5 Experimentations and further directions of research

Experimentations: counter-examples and patterns

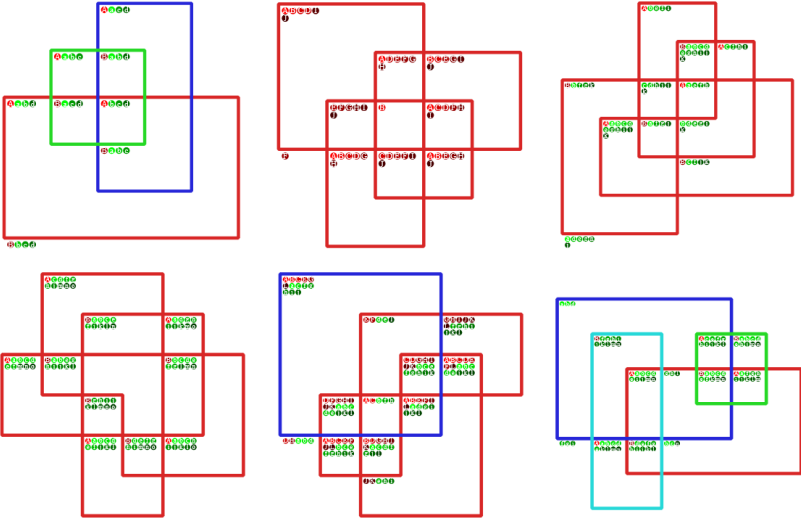
An online catalog: <https://christopherlloyd.github.io/LooPIndex/>

We constructed a data basis of $> 10^3$ multiloops with their pinning data. This enabled us to find counter-examples to various naïve conjectures:

- invariance under Reidemeister moves or mutations
- monotonicity properties of pinning numbers and their averages
- Some distinct loops have isomorphic pinning semi-lattices

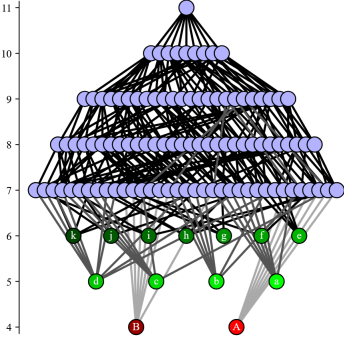
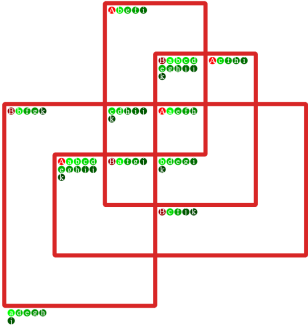
but many patterns are to be understood...

Experimentations: counter-examples and patterns



Multiloops with regions of degree > 2 ...

Experimentations: counter-examples and patterns



... have a rich pinning semi-lattice !

Further directions of research

Complexity of the pinning number on special instances

- 1 Find interesting classes of loops for which the LoopPinning problem is P.
(Add local or global restrictions of topological or geometrical nature.)
- 2 Generalise the MoBiDisc theorem to certain classes of multiloops.
(This requires extending some of the results by Hass-Scott.)

What do pinning quantities say about plane loops ?

- 1 Which are related to other invariants of plane loops ?
(Relations with Arnold's cohomological invariants J^- , J^+ , St ?)
- 2 How do they vary under Reidemeister moves ?
(Study the minimal pinning number over an $R3$ -equivalence classes.)

Applications to bounded cohomology of braid groups and $\text{Homeo}(\mathbb{S}^2)$

For $\phi \in \text{Homeo}(\mathbb{S}^2)$: pinning numbers of multiloops $\phi(\gamma) \cup \gamma$ can measure "braiding index" or "distance" to \mathcal{B}_p , translation length on curve graphs...

Bibliography



V. Arnold.

Topological invariants of plane curves and caustics.
AMS, 2000.



S. J. Blank.

Extending immersions of the circle.
Ph.D. dissertation, 1967.
Brandeis University, Waltham, MA.



Dennis Frisch.

Classification of immersions which are bounded by curves in surfaces.
Ph.D. dissertation, 2010.



Joel Hass and Peter Scott.

Intersections of curves on surfaces.
Israel J. Math., 51(1-2):90–120, 1985.



Joel Hass and Peter Scott.

Shortening curves on surfaces.
Topology, 33(1):25–43, 1994.



Max Neumann-Coto.

A characterization of shortest geodesics on surfaces.
Algebr. Geom. Topol., 1:349–368, 2001.



Valentin Poénaru.

Extension des immersions en codimension 1 (d'après Samuel Blank).
In *Séminaire Bourbaki*, Vol. 10, pages Exp. No. 342, 473–505. Soc. Math. France, Paris, 1995.



Christopher-Lloyd Simon and Ben Stucky.

Pin the loop taut: a one-player topological game, 2024.
Submitted for publication, [arxiv version](#).